

Internet Architectural Philosophy and the New Business Reality

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Topics

- ◆ architecture (as design philosophy)
- ◆ key decisions
- ◆ architecture (as reality)
- ◆ and then there is money

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Background

- ◆ multiple unrelated efforts (early to mid 1960's)
 - packet switching theory: (Kleinrock) 1961
 - day dreaming: (Licklider's Galactic Network) 1962
 - make use of remote expensive computers: (Roberts) 1964
 - survivable infrastructure for voice and data: (Baron) 1964
- ◆ ARPANET (late 1960's)
 - Roberts ARPANET paper 1967
 - RFP for "Interface Message Processor" won by BBN 1968
 - four ARPANET hosts by 1969
 - public demo and email in 1972

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Fundamental Goal of Internet Protocols

- ◆ multiplexed utilization of **existing** networks
 - different administrative boundaries
 - multiplexing via packets
 - networks interconnected with packet switches
 - called gateways (now called routers)
 - note: international in scope
- ◆ did not want to build a new global network
 - too expensive
 - too limiting

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Internet Protocols Design Philosophy

- ◆ ordered set of 2nd-level goals
 - 1/ **survivability** in the face of failure
 - 2/ support **multiple types** of communications service
 - 3/ accommodate a **variety** of network types
 - 4/ permit **distributed management** of resources
 - 5/ **cost effective**
 - 6/ **low effort** to attach a host
 - 7/ **account** for use of resources
- ◆ note: no performance (QoS) or security goals
- ◆ not all goals have been met
 - management & accounting functions are limited

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Packets!

- ◆ basic decision: use packets not circuits
 - Kleinrock's work showed packet switching to be a more efficient switching method
- ◆ packet (a.k.a. datagram)
 - self contained
 - handled independently of preceding or following packets
 - contains destination and source **internetwork** address
 - may** contain processing hints (e.g. QoS tag)
 - no delivery guarantees**
 - net may drop, duplicate, or deliver out of order
 - reliability (where needed) is done at higher levels

Dest Addr	Src Addr	payload
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Routing

- ◆ sub parts of the network are connected together by computers that forward packets toward destination these computers are called “**routers**”
- ◆ routers use destination address in packet to make forwarding decision
- ◆ routers exchange reachability information with other routers to build tables of “next hops” toward specific local networks
 - exchange of reachability information done with “**routing protocol**”

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Unreliability can be Important

- ◆ basic decision: **offer** an unreliable service
- ◆ 1st idea was to only have TCP (a reliable service)
- ◆ problems
 - not good for voice & video
 - data has to be delivered in time - retransmission for reliability causes too great a delay
 - not good for all applications
 - e.g. a debugger has to work in lossy environment
 - retransmission algorithm may vary with application
- ◆ thus: **split** IP & TCP and add UDP
 - IETF just added SCTP

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A Quote

*“the lesson of the Internet is that **efficiency is not the primary consideration**. Ability to grow and adapt to changing requirements is the primary consideration. This makes simplicity and uniformity very precious indeed.”*

Bob Braden

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Networks as Generic

- ◆ design requirement of working over:
existing networks & a wide variety of networks
- ◆ minimum set of assumptions about network
reasonable size packets, reasonable but not perfect
delivery reliability, network-wide addressing, way to
get error messages back to source, no assumption of in-order packet delivery
- ◆ “smart wires” are not much of a help
e.g. X.25 (reliable delivery)
e.g. ATM (QoS functions)
- ◆ thus it is easy to use new types of networks
assuming they are not too helpful (feature rich)

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End-to-End Argument

- ◆ 1981 paper by Saltzer, Reed & Clark
- ◆ “smart networks” do not help
 - adding functions into network can be redundant since actual function is end-to-end
 - e.g. encryption, data reliability
 - also harder to change with new technology
 - also see Lampson *Hints for Computer System Design*
- ◆ e2e argument projected to mean
 - no per-session knowledge or state in the network
 - but some “soft-state” (auto refreshed) may be OK
 - network should be transparent to end-to-end applications

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Ease of Experimentation With e2e

- ◆ easier to experiment in an e2e environment
 - if the network is transparent then only nodes involved are the end nodes
 - note that an end node could be a 3rd party server
 - no need to get permission to experiment
- ◆ cheaper to experiment
 - can do much smaller scale experiments - down to 2 nodes than core-based services
- ◆ WWW an example of what can be done

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Economic Driver?

- ◆ Mark Gaynor Harvard PHD thesis
- ◆ define market uncertainty as **MU**
 - how well do you know what the customer wants
- ◆ low MU means customer wants are known
 - e.g. “voice service”
 - no opportunity to be “better” than competitor
- ◆ high MU means customer wants are not known
 - e.g. future IP-enabled voice service
 - opportunity to better match customer wants than competitor does

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Economic Driver, contd.

- ◆ low MU
 - commodity service
 - provide most efficient way - frequently centralized
- ◆ high MU
 - need to experiment to try to match customer want
 - note: if only one company figures it out they dominate the market
 - easier to experiment on edges
 - i.e. e2e is a innovation friendly model
 - even if its more expensive to provide service to ends

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Smart vs. Stupid Networks

- ◆ phone network technology: self-named “Intelligent Network” (IN)
 - many network-based services
 - admission control, number translation, accounting, ...
- ◆ Isenberg’s *Rise of the Stupid Network* compared phone network’s “Intelligent Network” to Internet
 - Isenberg’s basic messages:
 - network (i.e. carrier) -based services slow to change
 - voice is not all there is
 - carrier gets in the way
 - just “deliver the bits” works

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But!!

- ◆ a “stupid network” is a commodity service
 - the price of a commodity service is driven by the stupidest vendor
- ◆ hard to make money delivering commodity services
- ◆ new network infrastructure is very expensive
 - fiber optic cables (with installation) & hardware
- ◆ access rights can also be very expensive
 - e.g. wireless spectrum licenses
- ◆ carriers need something else to make money
 - common dream is that services or content will save the day
 - may be a false dream



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But!! (2)

- ♦ packets w/o circuits cause problems
 - can not do guaranteed QoS
 - can not control path packets take
 - can not reserve capacity for application
 - security control harder
 - do not have logical “wire” back to source
 - management harder
 - can not see data patterns on the network
 - finding non-catastrophic failures harder
 - service provider interconnections harder
 - no clean interface
- ♦ lack of useful formal tools to describe performance

!QoS

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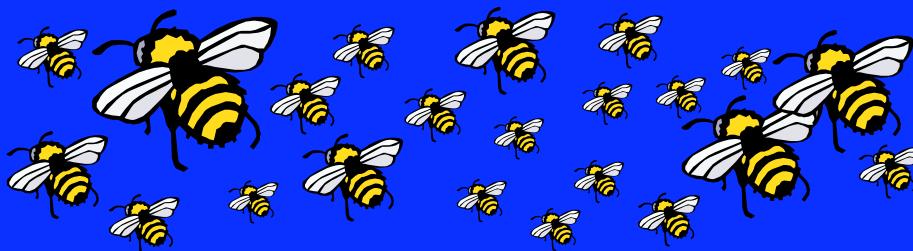
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Conceptualization Problem

- ♦ fundamental disconnect between “Internet” and “phone” people “bell-heads vs. net-heads”
- ♦ by their definition the Internet can not work and must be fixed - they will rescue us

“You can not build corporate network out of TCP/IP.”

IBM circa 1992

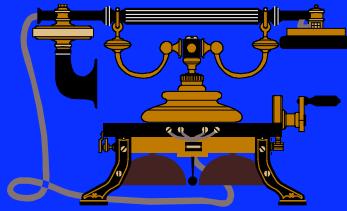


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Traditional Phone Network

- ♦ circuits & “smart network”
- ♦ connection-oriented
- ♦ hard state in network devices
- ♦ fragile
- ♦ central resource control
- ♦ socialist? "for the good of all"
- ♦ applications in network
 - e.g., phone switch
 - end-to-end touch-tone signaling was a mistake
- ♦ predictable development path
 - extended development cycle

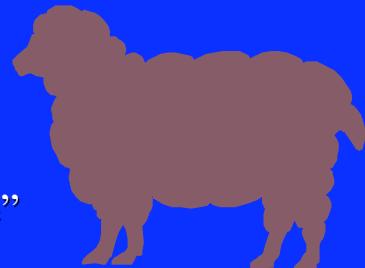


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Internet

- ♦ packets & e2e
- ♦ soft state in network devices
- ♦ resilient
- ♦ competitive resource control
- ♦ capitalist? "individual initiative"
 - but too much selfishness hurts all
 - must play by the same rules - but no enforcement
- the tragedy of the commons**
- ♦ applications in hosts at edges (end-to-end)
 - and in 3rd party servers anywhere on the net
- ♦ hard to predict developments
 - chaos at the rate of “Internet time”



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More Conceptualization Problems

- ◆ service provided by 3rd parties - not only by carriers
 - different from phone world
- ◆ a quote from an IETF telephony mailing list

Hi Roy,
I still don't understand why it is a "users" choice where the "services" are executed -
I would have thought that this would be networks choice

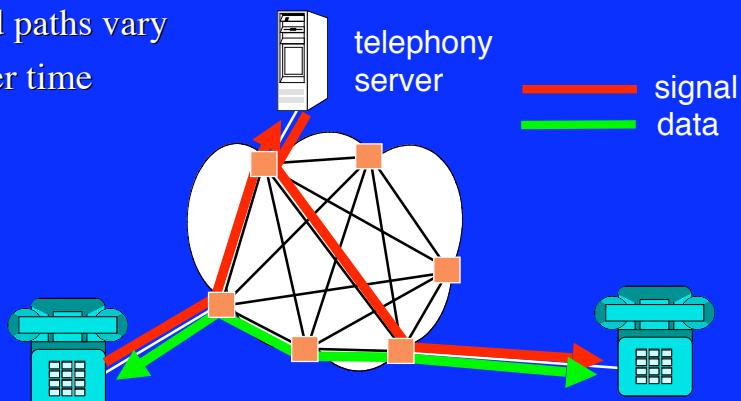
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Disjoint Control and Data Paths

- ◆ signaling and data paths in Internet may not coincide

and paths vary
over time



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Circuits in the Internet

- ♦ do not seem to go away (MPLS)
- ♦ used for traffic engineering
 - city-pair pipes
 - maybe class of service city-pair pipes
- ♦ and customer connections
- ♦ finer grain (instance of application) use still pushed
- ♦ remember the fate of ATM
 - circuit - used for trunks not flows
 - QoS - ignored (ATM not end-to-end)
 - link sharing - may make sense
 - as the bearer service - did not make it



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IP as a Common Bearer Service

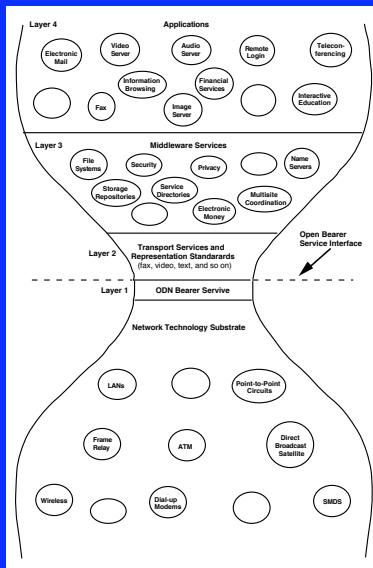


FIGURE 2.1 A four-layer model for the Open Data Network

From: Realizing the
Information Future

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Net is No Longer Transparent

- ◆ end-to-end argument says the net should be transparent
 - i.e. packet not modified in transit (other than TTL)
 - global-scope internetwork address
 - i.e., packet goes to address in destination address field
- ◆ transparency now gone in some cases
 - NATs, firewalls, proxies, content caches, TCP reshapers replace addresses, intercept traffic, insert traffic
- ◆ other issues
 - wiretapping, taxation, content filtering

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NAT/Firewall/Cache Issues

- ◆ can not trust IP address as end-to-end
 - breaks IPSec, not sure who you are talking to
- ◆ applications with addresses in data
 - have to have application-specific support (ALG) in devices
 - deploying new application requires approval of net manager
- ◆ dynamic port usage
 - ALG must snoop on application traffic
 - ALG must understand application logic
- ◆ new IETF effort to develop generic signaling
 - may help some
 - but will not make these devices transparent

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Trust-Free Environment

- ♦ original Internet architecture assumed a trustworthy environment
- ♦ no longer the case
 - mistrust net itself (eavesdropping, reliability etc)
 - mistrust that you are talking to the right end point
 - e.g., proxy, redirect, spoofing (MAC & IP address)
 - unsolicited correspondence (spam)
 - anonymity hard to get
 - mistrust own hardware and software
 - 3rd parties insist on being in the middle
 - filters, wiretapping, ...

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Summary of Architectural Points

- ♦ datagram-based network
 - not circuit switched
- ♦ network of networks
 - different parts under different management
- ♦ minimize per-session state in network
 - some auto-refreshed state is OK
- ♦ end-to-end model maximizes flexibility
 - network does not need to know what you are doing
- ♦ “smart wires” can get in the way
 - e.g., nested control loops
- ♦ reliable delivery is an option
 - not a requirement

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Key Decisions

- ◆ a few key decisions brought us here to the Internet of today
- ◆ but there was no way to predict where we are now
- ◆ **unplanned parenthood**

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10 Decisions That Made a Difference

- ◆ support existing networks
- ◆ **datagram**-based
- ◆ creating the **router** function
- ◆ split TCP **and** IP
- ◆ DARPA fund Berkeley to add TCP/IP to **UNIX**
- ◆ CSNET and **CSNET/ARPANET** deal
- ◆ NSF **require TCP/IP** on NSFnet
- ◆ ISO **turn down** TCP/IP
- ◆ NSF Acceptable Use Policy (**AUP**)
- ◆ **minimal** regulation

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Internet Architecture

- ◆ #1 goal of original Internet protocols was to deal with a **network of networks**
 - not a single type of network
 - not under one management
- ◆ networks interconnected at datagram level
 - no session-aware logic at interconnections
- ◆ bi-lateral interconnection agreements
 - “customer” - buy transit service to “the Internet”
 - “peer” - cost sharing connection to a network and its customers

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Customer Interconnection

- ◆ one network pays another for access to “the Internet”
 - paying network can be Internet service provider (ISP) or enterprise
 - only as useful as resulting coverage
 - “Metcafe’s Law”: value of network increases by square of the number of reachable nodes
- ◆ customer can move business to another network if they do not like the service
 - may have to renumber to preserve addressing topology

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Peering

- ◆ business decision
 - no current regulations
- ◆ it can be cost effective for two networks to interconnect sharing the costs of the links
 - interconnection can be at “public peering points” or using dedicated links between networks
- ◆ but only “see” other network and their customers
 - not the other network’s other peers
- ◆ must peer with all large networks to get “the Internet”
 - or be a customer to another network (or networks)

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Multi-Homing

- ◆ one network (ISP or enterprise) can connect to more than one other network
 - for redundancy and reliability
 - called “multi-homed”
- ◆ causes some complexity in the routing setup

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Public Peering Points

- ◆ 3 originally designated by National Science Foundation (NSF) as part of the breakup of the NSFnet
- ◆ now many local peering points around the world
 - but telcom costs can discourage use in some countries
 - cheaper to get lines to US than within country
- ◆ level-2 interconnect
 - like an local area network (e.g. an Ethernet)
 - i.e. not involved in IP-level routing

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Private Peering

- ◆ two ISPs can agree to interconnect sharing costs
 - “you buy and run one line, I’ll buy and run another”
 - peering list normally private
- ◆ ISPs have minimum criteria before peering will be considered
 - some publish the criteria
- ◆ criteria normally include
 - minimum level of interconnect traffic, traffic balance,
 - backbone size, geographic scope,
 - competent network operations center

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Tier 1 ISPs

- ◆ some big ISPs are referred to as “Tier 1 ISPs”
- ◆ no real externally verifiable definition
- ◆ general concept:
 - “an ISP that gets most of not all of its connectivity from peering, not by being a transit customer”
 - i.e. a Tier 1 ISP is one that is connected to the other Tier 1 ISPs

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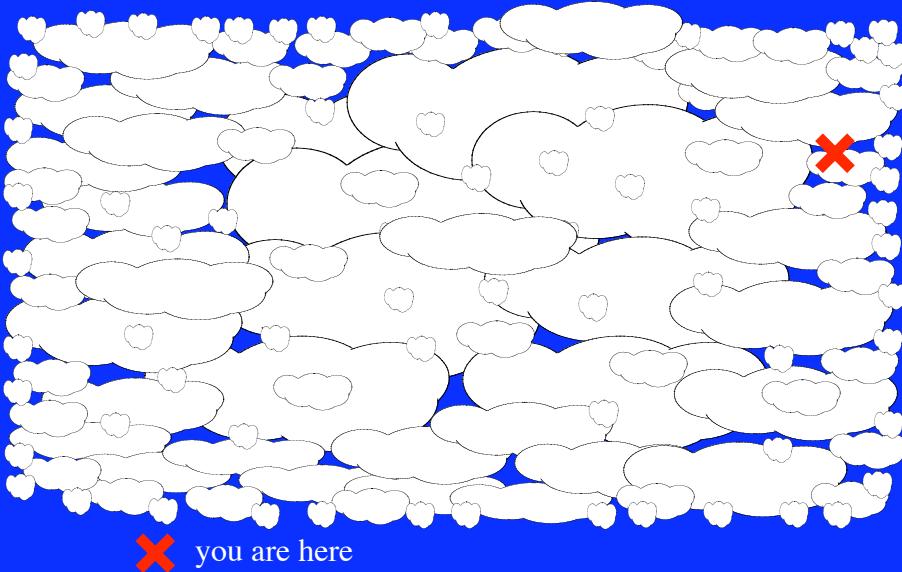
Interconnection Pattern

- ◆ no explicit network hierarchy assumed
- ◆ no specific pattern to ISP interconnections
 - other than that peering tends to be between networks of the same basic size
 - but not always - can have business reasons for mismatch
- ◆ peering and transit connections can appear random
- ◆ **notes:**
 - most traffic does not flow through Tier 1 ISPs
 - many “lower-level” interconnections
 - hard (impossible) to know relative sizes of ISPs

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Current Internet Architecture



Money

- ♦ “but who is going to make money at that?”

John McQuillan

- ♦ how is the carrier supported?

“we do not know how to route money”

Dave Clark

- ♦ carrier wants a piece of the action

e.g., WAP, AT&T proposal

- ♦ is content king?

factoid: total US movie revenue \approx 2 weeks of US phone charges

iMode: A Model?

- ◆ DoCoMo's iMode service
 - more than 30 million subscribers
 - 9.6Kb data service
 - 50,000 iMode compatible sites
 - DoCoMo works with less than 10% of them
 - does billing, runs servers etc
 - rest are on their own
- ◆ key decision: open access (NOT WAP!)
 - makes service more attractive
 - DoCoMo charges monthly fee and for data transferred

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More on Money

- ◆ QoS does not seem to be a useful charging base
- ◆ differentiated by application is an intelligence test
 - railroads in US used to do this (Rhode Island Line )
 - not enough will fail the test
- ◆ and then there is all that fiber
- ◆ do municipalities have a role?

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A Bit More on Money

- ◆ what happened to that \$ trillion anyway?
 - few infrastructures pay for themselves
 - the Internet is not an exception
- ◆ is there a difference now that the fiber is “free”?

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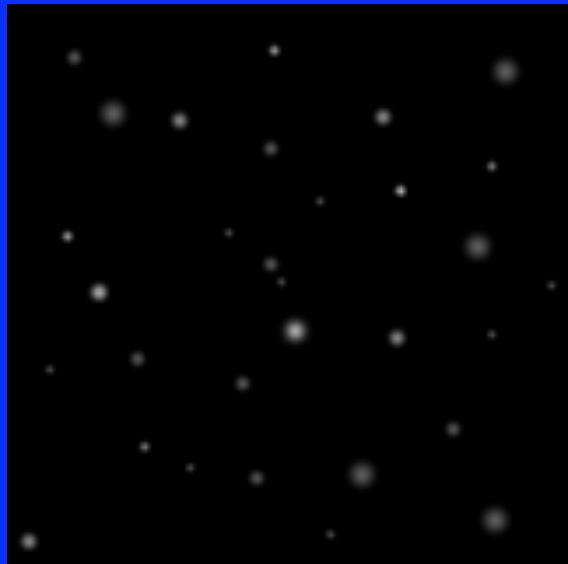
Last Word

- ◆ Internet “too important to fail” (?)
 - what about ISPs (can you say “KPNQuest”?)
- ◆ will there be anyone left standing other than the telcos?
 - what can they see from their point of view?
- ◆ will you be able to say “Internet” and “business model” in the same sentence?
 - without a “no” in between

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What's Next?



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